

## 8.96: Grover's Search Algorithm- Four-Card Monte

Grover's search algorithm is great at playing four-card monte. As the following quantum circuit shows it can determine which card is the queen in one pass.

$$\begin{array}{l}
 |0\rangle \xrightarrow{H} \left[ \begin{array}{c} \phantom{0} \\ \text{Oracle} \end{array} \right] \xrightarrow{H} \left[ \begin{array}{c} \phantom{0} \\ J \end{array} \right] \xrightarrow{H} \text{Measure} \quad \left[ \begin{array}{c} \phantom{0} \\ \phantom{0} \end{array} \right] \quad X \cdot X \\
 |0\rangle \xrightarrow{H} \left[ \begin{array}{c} \phantom{0} \\ \phantom{0} \end{array} \right] \xrightarrow{H} \left[ \begin{array}{c} \phantom{0} \\ \phantom{0} \end{array} \right] \xrightarrow{H} \text{Measure} \quad \left[ \begin{array}{c} \phantom{0} \\ \phantom{0} \end{array} \right] \quad X \boxed{Z} X
 \end{array}$$

where  $\left[ \begin{array}{c} \phantom{0} \\ J \end{array} \right] = \left[ \begin{array}{c} \phantom{0} \\ \phantom{0} \end{array} \right]$

The following matrix operators are required to construct the circuit. Giving  $|10\rangle$  a negative phase in the Oracle designates it as the queen. The Appendix shows the calculation of J as shown on the right side above.

$$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix} \quad \text{HH} = \text{kroncker}(H, H) \quad \text{Oracle} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & -1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \quad J = \begin{pmatrix} -1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

Operating on the input state, which creates a superposition of all queries, enables the algorithm to identify which card is the queen in one operation of the circuit.

$$\text{GroverSearch} = \text{HH} J \text{HH Oracle HH} \quad \text{GroverSearch} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ -1 \\ 0 \end{pmatrix} = -|10\rangle$$

Now the operation of the algorithm is carried out in stages to show the importance of constructive and destructive interference in quantum computers.

$$\begin{array}{l}
 \text{Step 1} \quad \text{HH} \begin{pmatrix} 1 \\ 0 \\ 0 \\ 0 \end{pmatrix} = \begin{pmatrix} 0.5 \\ 0.5 \\ 0.5 \\ 0.5 \end{pmatrix} = \frac{1}{2}[|00\rangle + |01\rangle + |10\rangle + |11\rangle] \\
 \text{Step 2} \quad \text{Oracle} \begin{pmatrix} 0.5 \\ 0.5 \\ 0.5 \\ 0.5 \end{pmatrix} = \begin{pmatrix} 0.5 \\ 0.5 \\ -0.5 \\ 0.5 \end{pmatrix} = \frac{1}{2}[|00\rangle + |01\rangle - |10\rangle + |11\rangle] \\
 \text{Step 3} \quad \text{HH} \begin{pmatrix} 0.5 \\ 0.5 \\ -0.5 \\ 0.5 \end{pmatrix} = \begin{pmatrix} 0.5 \\ -0.5 \\ 0.5 \\ 0.5 \end{pmatrix} = \frac{1}{2}[|00\rangle - |01\rangle + |10\rangle + |11\rangle] \\
 \text{Step 4} \quad J \begin{pmatrix} 0.5 \\ -0.5 \\ 0.5 \\ 0.5 \end{pmatrix} = \begin{pmatrix} -0.5 \\ -0.5 \\ 0.5 \\ 0.5 \end{pmatrix} = \frac{1}{2}[-|00\rangle - |01\rangle + |10\rangle + |11\rangle] \\
 \text{Step 5} \quad \text{HH} \begin{pmatrix} -0.5 \\ -0.5 \\ 0.5 \\ 0.5 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ -1 \\ 0 \end{pmatrix} = -|10\rangle
 \end{array}$$

### Appendix

$$\left[ \begin{array}{c} \phantom{0} \\ J \end{array} \right] = \left[ \begin{array}{c} \phantom{0} \\ \phantom{0} \end{array} \right]$$

$$J = \begin{pmatrix} -1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \quad X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad CZ = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -1 \end{pmatrix} \quad \text{kronecker}(X, X) CZ \text{kronecker}(X, X) = \begin{pmatrix} -1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix}$$

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